

Discrete Mathematics

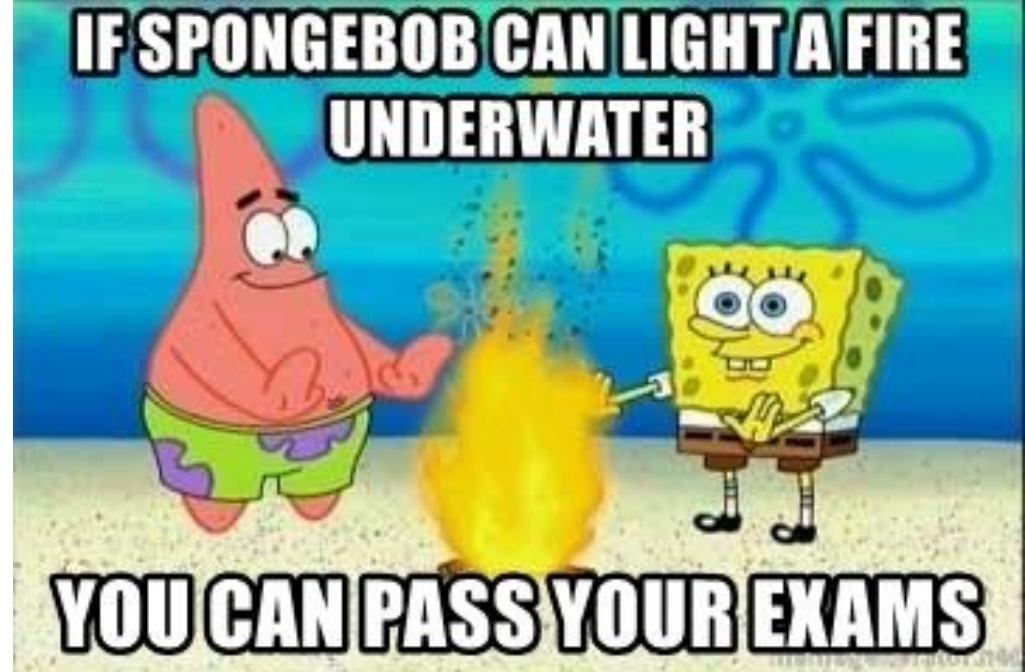
(Final Session)

Zhaoguo Wang

范围1: 期末考试 (闭卷)

- 逻辑1, 2, 4, 5章
 - 不包括存在型前束范式
- 集合论9, 10, 11.1, 11.2章节

- 一切以课本为准, 不要使用switch variables等课本上没有的技术



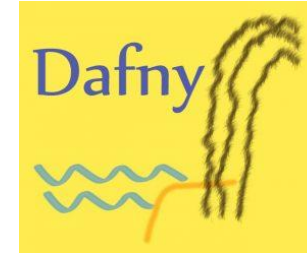
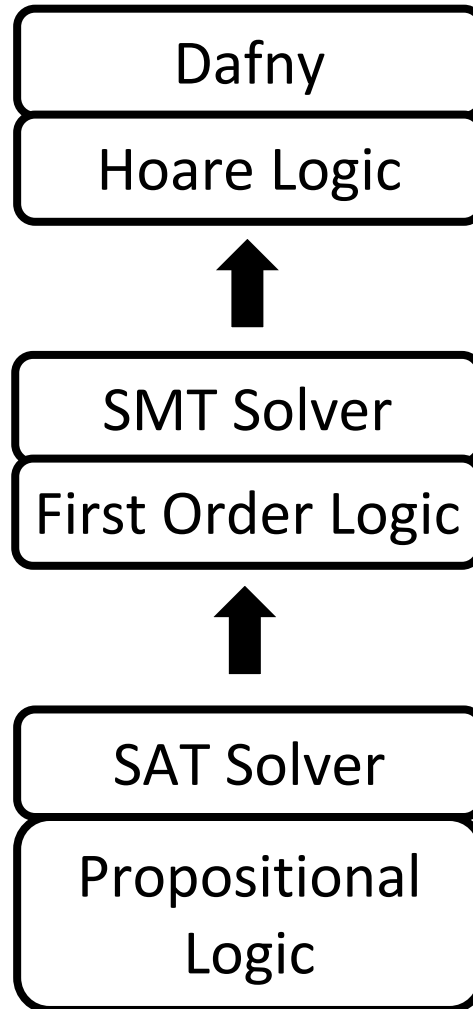
范围2: 测验 (开卷)

- 形式化验证
 - DPLL, convert program into SAT
 - Eager SMT, Lazy SMT, EUF, symbolic execution,
 - Hoare triple, Hoare inference rules, loop invariant(quiz会给出invariant)
 - 不考 SMT Quantifier, SMT trigger, weakest precondition, switch variables
- 自动机和图灵机
 - DFA, Single-tape Turing Machine
 - 不考 Turing Machine Variants



Verification Architecture

Z3



The
MiniSat
Page

Part I: Computer Aided Reasoning Overview

1.1 Propositional Logic

Step 1. Convert program into mathematical formula.

Step 2. Ask the computer to solve the formula.

1.2 First Order Logic

Step 1. Convert program into first order logic formula.

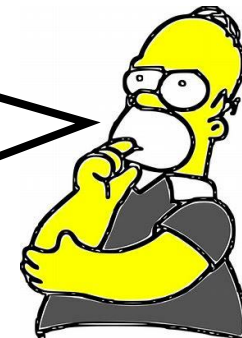
Step 2. Ask the computer to solve the formula.

1.3 Auto-active Proof

Step 1. Axiom system

Step 2. Ask the computer to check the invariants

We start at propositional logic, which is a simple logic. Computer has 0 and 1. Propositional logic has T and F.



Propositional Logic

- A **proposition(命题)** is a statement that is, by itself, either true or false.
- Every formula in propositional logic consists of **propositional variables (命题变项)** combined via **propositional connectives (联结词)**.
 - $p, q, \dots$
 - $\neg, \wedge, \vee, \rightarrow, \leftrightarrow$
- Propositions without any connectives are called **atomic propositions or simple propositions(原子命题/简单命题)**.

Well-Formed Formulas (合式公式)

INDUCTIVE DEFINITION of WFF

- 1). Every single proposition (symbol) is WFF.
- 2). If A and B are WFF, so are $(\neg A)$, $(A \wedge B)$, $(A \vee B)$, $(A \rightarrow B)$, and $(A \leftrightarrow B)$.
- 3). No expression is WFF unless forced by 1) or 2).

- We can omit some brackets.
- All operators are left-associative.
- Operator precedence: $\neg > \wedge > \vee > \rightarrow > \leftrightarrow$

Truth Table (真值表)

- The semantic of a WFF can be represented as a truth table.
- If there are n propositional variables, the truth table has 2^n rows.
- Truth tables and formulas can be converted to each other.

P	Q	$g_0(P, Q)$	$g_1(P, Q)$
F	F	T	T
F	T	T	T
T	F	F	F
T	T	T	F

$$g_0(P, Q) = (((\neg P) \wedge (\neg Q)) \vee ((\neg P) \wedge Q)) \vee (P \wedge Q)$$

$$g_1(P, Q) = ((\neg P) \wedge (\neg Q)) \vee ((\neg P) \wedge Q)$$

M1. According to the T rows

P	Q	$g_0(P, Q)$	$g_1(P, Q)$
F	F	T	T
F	T	T	T
T	F	F	F
T	T	T	F

$$g_0(P, Q) = ((\neg P) \vee Q)$$

$$g_1(P, Q) = ((\neg P) \vee Q) \wedge ((\neg P) \vee (\neg Q))$$

M2. According to the F rows

Truth Table

- We can also define new connectives.
- For n propositional variables, we can define 2^{2^n} connectives.
- But some connectives are unnecessary.

P	Q	$P \nabla Q$
T	T	F
F	T	T
T	F	T
F	F	F

$$P \nabla Q = ((P \wedge (\neg Q)) \vee ((\neg P) \wedge Q))$$

Completeness of Connectives (完备性)

DEFINITION

A group of connectives are **complete** if

- 1). Every formula can be converted to an **equivalent** formula and
- 2). The formula only includes connectives in the group.

- $\{\neg, \wedge\}$ and $\{\neg, \vee\}$ are all complete.
- More examples in textbook 2.4.
- We can determine equivalence via **truth table** or **calculation via laws**.

Laws

De Morgan's Law (摩根律):

$$\neg(P \wedge Q) = \neg P \vee \neg Q, \neg(P \vee Q) = \neg P \wedge \neg Q$$

Distributive Law (分配律) :

$$P \vee (Q \wedge R) = (P \vee Q) \wedge (P \vee R)$$

$$P \wedge (Q \vee R) = (P \wedge Q) \vee (P \wedge R)$$

Absorption Law (吸收律) :

$$P \vee (P \wedge Q) = P$$

$$P \wedge (P \vee Q) = P$$

More laws in 2.2

Laws

$$\begin{aligned} & P3 \wedge \neg(P1 \wedge P2) \wedge \neg(P2 \wedge P3) \wedge \neg(P1 \wedge P3) \wedge \neg(\neg P1 \wedge \neg P2) \\ = & P3 \wedge (\neg P1 \vee \neg P2) \wedge (\neg P2 \vee \neg P3) \wedge (\neg P1 \vee \neg P3) \wedge (P1 \vee P2) \\ = & P3 \wedge (\neg P2 \vee \neg P3) \wedge (\neg P1 \vee \neg P2) \wedge (\neg P1 \vee \neg P3) \wedge (P1 \vee P2) \\ = & P3 \wedge \neg P2 \wedge (\neg P1 \vee \neg P2) \wedge (\neg P1 \vee \neg P3) \wedge (P1 \vee P2) \\ = & P3 \wedge \neg P2 \wedge (P1 \vee P2) \wedge (\neg P1 \vee \neg P2) \wedge (\neg P1 \vee \neg P3) \\ = & P3 \wedge \neg P2 \wedge P1 \wedge (\neg P1 \vee \neg P2) \wedge (\neg P1 \vee \neg P3) \\ = & P3 \wedge \neg P2 \wedge P1 \wedge (\neg P1 \vee \neg P3) \wedge (\neg P1 \vee \neg P2) \\ = & P3 \wedge \neg P2 \wedge P1 \wedge \neg P3 \wedge (\neg P1 \vee \neg P2) \\ = & P3 \wedge \neg P3 \wedge \neg P2 \wedge P1 \wedge (\neg P1 \vee \neg P2) \\ = & F \end{aligned}$$

It is a special proposition
which is always false

How to Prove

DEFINITION

Tautology (重言式/永真式) is a proposition which is always true under any interpretation. $P \vee \neg P$

Contradiction (矛盾式) is a proposition which is always false under any interpretation. $P \wedge \neg P$

If a proposition can be true under some interpretation, it is **satisfiable (可满足)**.

$$P3 \wedge \neg(P1 \wedge P2) \wedge \neg(\neg P1 \wedge P3)$$

- A is a tautology iff. $\neg A$ is a contradiction (**unsatisfiable**).
- We can represent a problem as a WFF and prove it by
 - Truth table (very complex because of **2^n** rows)
 - SAT solver
 - Deduction

SAT Solver

DEFINITION

In computer science, the satisfiability problem (abbreviated SATISFIABILITY or SAT) is the problem of determining if there exists an interpretation that satisfies a given formula.

- SAT solver is used to solve SAT problem.
- Given a WFF, SAT solver returns “sat” or “unsat”.
- The WFF must be in a normal form (范式).

Normal Form

- A literal (文字) is a propositional variable or its negation.
- A clause (析取式/子句) is a disjunction of one or more literals.
- A conjunctive clause (合取式) is a conjunction of one or more literals.
- Conjunctive normal form (合取范式) $(A \vee \neg B) \wedge (B \vee \neg A) \wedge A$
 - A conjunction of one or more clauses.
- Disjunctive normal form (析取范式) $(A \wedge B) \vee (\neg A \wedge \neg B)$
 - A disjunction of one or more conjunction clauses.
- Every formula can be converted to CNF or DNF.
- SAT solver uses CNF because
 - Converting a formula to DNF may result in a much larger formula.

Principle Normal Form (主范式)

- Given n propositional variables, if every variable exists once in a conjunction clause, the clause is a **minterm (极小项)**
- Principal Disjunction Normal Form (主析取范式)
 - **Disjunction of minterms** $g_0(P, Q) = (((\neg P) \wedge (\neg Q)) \vee ((\neg P) \wedge Q)) \vee (P \wedge Q)$
- Given n propositional variables, if every variable exists once in a disjunction clause, the clause is a **maxterm (极大项)**
- Principal Conjunction Normal Form (主合取范式)
 - **Conjunction of maxterms** $g_0(P, Q) = ((\neg P) \vee Q)$
- Every formula with n propositional variables has one and only one PCNF and PDNF
- SAT solver do not use PNF because CNF is enough

DPLL

DEFINITION

A literal is defined if its truth value has been guessed or inferred.

OVERVIEW

Iteratively process by following steps:

- Guess the truth value of an undefined literal
- Infer the truth value of other variables with inference rules
- If a clause is false and there is a decision literal, backtrack and re-guess a variable

DPLL

$$\emptyset \parallel 1 \vee 3, \bar{2} \vee \bar{4}, \bar{1} \vee 2 \vee \bar{4}, \bar{1} \vee 3, \bar{2} \vee \bar{3} \vee 4, 1 \vee 4, \bar{1} \vee 2 \vee 4 \Rightarrow (\textit{Decide})$$

$$1^d \parallel 1 \vee 3, \bar{2} \vee \bar{4}, \bar{1} \vee 2 \vee \bar{4}, \bar{1} \vee 3, \bar{2} \vee \bar{3} \vee 4, 1 \vee 4, \bar{1} \vee 2 \vee 4 \Rightarrow (\textit{UnitPropagate})$$

$$1^d 3 \parallel 1 \vee 3, \bar{2} \vee \bar{4}, \bar{1} \vee 2 \vee \bar{4}, \bar{1} \vee 3, \bar{2} \vee \bar{3} \vee 4, 1 \vee 4, \bar{1} \vee 2 \vee 4 \Rightarrow (\textit{Decide})$$

$$1^d 3 2^d \parallel 1 \vee 3, \bar{2} \vee \bar{4}, \bar{1} \vee 2 \vee \bar{4}, \bar{1} \vee 3, \bar{2} \vee \bar{3} \vee 4, 1 \vee 4, \bar{1} \vee 2 \vee 4 \Rightarrow (\textit{UnitPropagate})$$

$$1^d 3 2^d \bar{4} \parallel 1 \vee 3, \bar{2} \vee \bar{4}, \bar{1} \vee 2 \vee \bar{4}, \bar{1} \vee 3, \bar{2} \vee \bar{3} \vee 4, 1 \vee 4, \bar{1} \vee 2 \vee 4 \Rightarrow (\textit{BackTrack})$$

$$1^d 3 \bar{2} \parallel 1 \vee 3, \bar{2} \vee \bar{4}, \bar{1} \vee 2 \vee \bar{4}, \bar{1} \vee 3, \bar{2} \vee \bar{3} \vee 4, 1 \vee 4, \bar{1} \vee 2 \vee 4 \Rightarrow (\textit{UnitPropagate})$$

$$1^d 3 \bar{2} \bar{4} \parallel 1 \vee 3, \bar{2} \vee \bar{4}, \bar{1} \vee 2 \vee \bar{4}, \bar{1} \vee 3, \bar{2} \vee \bar{3} \vee 4, 1 \vee 4, \bar{1} \vee 2 \vee 4 \Rightarrow (\textit{BackTrack})$$

$$\bar{1} \parallel 1 \vee 3, \bar{2} \vee \bar{4}, \bar{1} \vee 2 \vee \bar{4}, \bar{1} \vee 3, \bar{2} \vee \bar{3} \vee 4, 1 \vee 4, \bar{1} \vee 2 \vee 4 \Rightarrow (\textit{UnitPropagate})$$

$$\bar{1} 3 \parallel 1 \vee 3, \bar{2} \vee \bar{4}, \bar{1} \vee 2 \vee \bar{4}, \bar{1} \vee 3, \bar{2} \vee \bar{3} \vee 4, 1 \vee 4, \bar{1} \vee 2 \vee 4 \Rightarrow (\textit{UnitPropagate})$$

$$\bar{1} 3 4 \parallel 1 \vee 3, \bar{2} \vee \bar{4}, \bar{1} \vee 2 \vee \bar{4}, \bar{1} \vee 3, \bar{2} \vee \bar{3} \vee 4, 1 \vee 4, \bar{1} \vee 2 \vee 4 \Rightarrow (\textit{UnitPropagate})$$

$$\bar{1} 3 4 \bar{2} \parallel 1 \vee 3, \bar{2} \vee \bar{4}, \bar{1} \vee 2 \vee \bar{4}, \bar{1} \vee 3, \bar{2} \vee \bar{3} \vee 4, 1 \vee 4, \bar{1} \vee 2 \vee 4$$

Deduction (推理演算)

Prove R when $P \rightarrow Q$, $Q \rightarrow R$ and P .

(1) P

introduce premise

前提引入

(2) $P \rightarrow Q$

introduce premise

分离规则

(3) Q

Hypothetical reasoning on (1)(2)

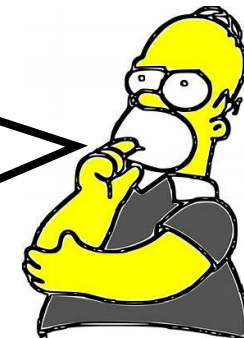
(4) $Q \rightarrow R$

introduce premise

(5) R

Hypothetical reasoning on (3)(4)

We can also prove a formula by deduction.



Deduction (归结推理法)

For formula $A = P \vee Q$ and $B = \neg P \vee R$, $R(A, B) = Q \vee R$ is a resolvent(归结式) of A and B .

To prove $C \Rightarrow D$, we just need to prove $C \wedge \neg D$ is a contradiction.

Prove that $(P \rightarrow Q) \wedge P \Rightarrow Q$.

① $(P \rightarrow Q) \wedge P \wedge \neg Q = (\neg P \vee Q) \wedge P \wedge \neg Q$

② $S = \{\neg P \vee Q, P, \neg Q\}$

③

(1) $\neg P \vee Q$

introduce premise

(2) P

introduce premise

(3) Q

resolution of (1) and (2)

(4) $\neg Q$

introduce premise

(5) $\square$

resolution of (3) and (4)

Convert Program to WFF

```
void swap(bool& a, bool& b)
{
    a = a ^ b;
    b = b ^ a;
    a = a ^ b;
}
```

a0

b0

state0

$a = a \wedge b;$

a1

b1

state1

$b = b \wedge a;$

a2

b2

state2

$a = a \wedge b;$

a3

b3

state3

$$\begin{aligned} & (A1 \leftrightarrow (A0 \wedge \neg B0) \vee (\neg A0 \wedge B0)) \wedge \\ & (B1 \leftrightarrow B0) \wedge \\ & (A2 \leftrightarrow A1) \wedge \\ & (B2 \leftrightarrow (A1 \wedge \neg B1) \vee (\neg A1 \wedge B1)) \wedge \\ & (A3 \leftrightarrow (A2 \wedge \neg B2) \vee (\neg A2 \wedge B2)) \wedge \\ & (B3 \leftrightarrow B2) \rightarrow \\ & (A3 \leftrightarrow B0) \wedge (A0 \leftrightarrow B3) \end{aligned}$$

Program

Assertion

*a and b are swapped iff.
the WFF is a tautology.*

Dual Formulas (对偶式)

DEFINITION

Given a formula A , replace $\vee, \wedge, T, F$ in A with $\wedge, \vee, F, T$ and get A^* . Then A and A^* are dual formulas(对偶式) with each other.

THEOREM

$$A = A(P_1, \dots, P_n), \quad A^- = A(\neg P_1, \dots, \neg P_n)$$

$$\neg(A^*) = (\neg A)^*, \quad \neg(A^-) = (\neg A)^-$$

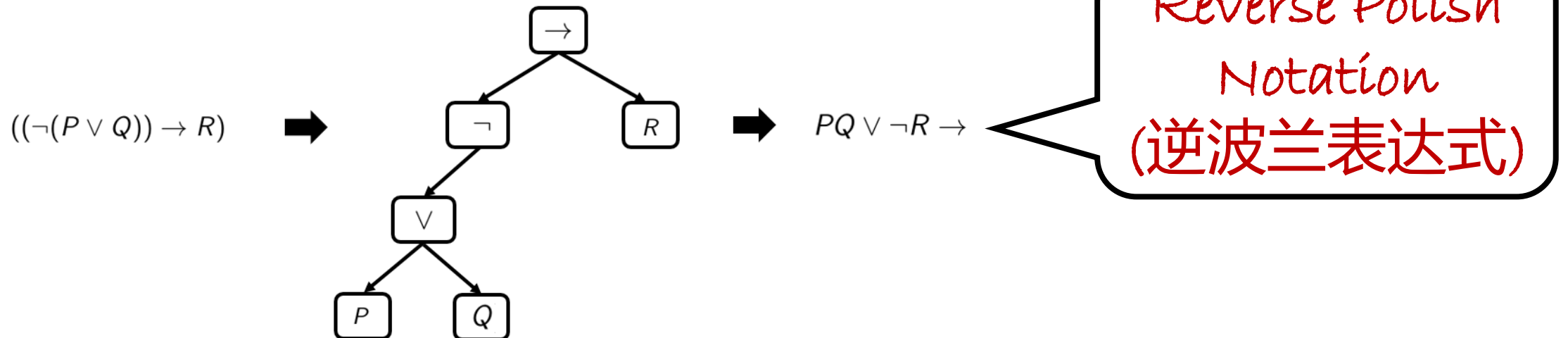
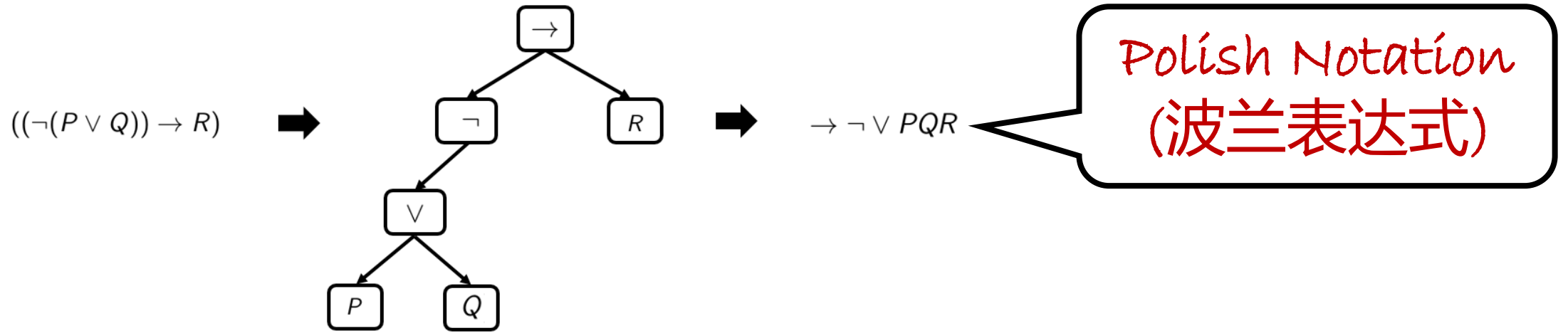
$$(A^*)^* = A, \quad (A^-)^- = A$$

$$(\neg A) = A^{*-}$$

More theorems in 2.5.



Polish Notation and Reverse Polish Notation



Part I: Computer Aided Reasoning Overview

1.1 Propositional Logic

Step 1. Convert program into mathematical formula.

Step 2. Ask the computer to solve the formula.

1.2 First Order Logic

Step 1. Convert program into first order logic formula.

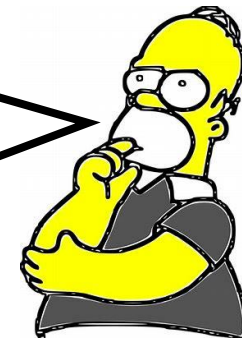
Step 2. Ask the computer to solve the formula.

1.3 Auto-active Proof

Step 1. Axiom system

Step 2. Ask the computer to check the invariants

Propositional logic cannot express complex statements conveniently. So we introduce predicate logic.



Predicate Logic

DEFINITION

Predicates (谓词) describe properties of individuals (个体词).

The range of individuals is domain of discourse (论域).

A function (函数) maps individuals to an individual instead of T or F.

	arguments	results
connective	propositions	a proposition
predicate	individuals	a proposition
function	individuals	an individual

Quantifier (量词)

DEFINITION

A quantifier turns a formula about individuals having some property into a formula about the number (quantity) of individuals having the property.

The blue rectangle is the scope (辖域) of "exists x ".

$(\exists x)((\forall y)(x \leq y))$

x and y are all bound variables (约束变元).

The red rectangle is the scope of "for all y ".

Well-formed Formula (合式公式)

INDUCTIVE DEFINITION of WFF

- 1). Every atomic formula is in WFF.
- 2). If A and B are WFFs, so are $(\neg A)$, $(A \wedge B)$, $(A \vee B)$, $(A \rightarrow B)$, and $(A \leftrightarrow B)$.
There is no variable which is bounded in one wff and free in the other wff.
- 3). If A is a WFF and x is free in A , then $(\forall x)A$, $(\exists x)A$ are wffs.
- 4). No expression is WFF unless forced by 1), 2) or 3).

Decision Problem (判定问题)

DEFINITION

If a formula is true under some interpretation, it is **satisfiable** (可满足).

If a formula is always false under any interpretations, it is **unsatisfiable** (不可满足).

If a formula is always true under any interpretations, it is **universally valid** (普遍有效).

Decision problem in predicate logic is whether there is an **effective algorithm** to determine if a formula is **universally valid**.

- Propositional logic is decidable but predicate logic is not decidable
- We can represent a problem as a WFF and prove it by
 - Deduction
 - SMT solver

Deduction (推理规则法)

Premise : $(\forall x)(P(x) \rightarrow Q(x)), (\forall x)(Q(x) \rightarrow R(x))$

Conclusion : $(\forall x)(P(x) \rightarrow R(x))$

Proof.

(1) $(\forall x)(P(x) \rightarrow Q(x))$

introduce premise

(2) $P(x) \rightarrow Q(x)$

remove universal quantifier

全称量词消去规则

(3) $(\forall x)(Q(x) \rightarrow R(x))$

introduce premise

(4) $Q(x) \rightarrow R(x)$

remove universal quantifier

三段论

(5) $P(x) \rightarrow R(x)$

sylogism on (2)(4)

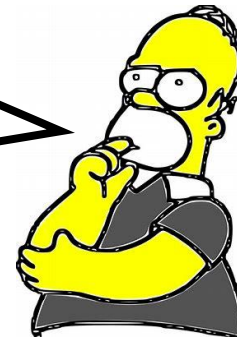
(6) $(\forall x)(P(x) \rightarrow R(x))$

introduce universal quantifier

全称量词引入规则

Deduction (归结推理法)

The basic idea is similar to that in propositional logic. But we should eliminate quantifiers firstly.



Prenex Normal Form (前束范式)

DEFINITION

A formula of the predicate calculus is in prenex normal form(PNF) if it is written as a string of quantifiers and bound variables, followed by a quantifier-free part, called the matrix (母式/基式).

quantifiers

$$(Q_1x_1) \dots (Q_nx_n)M(x_1, \dots, x_n)$$

Every formula can be converted into equivalent PNFs. (定理5.3.1)



Equivalence (等值)

$$(\forall x)(P(x) \vee q) = (\forall x)P(x) \vee q$$

$$(\exists x)(P(x) \vee q) = (\exists x)P(x) \vee q$$

$$(\forall x)(P(x) \wedge q) = (\forall x)P(x) \wedge q$$

$$(\exists x)(P(x) \wedge q) = (\exists x)P(x) \wedge q$$

$$(\forall x)(P(x) \rightarrow q) = (\exists x)P(x) \rightarrow q$$

$$(\exists x)(P(x) \rightarrow q) = (\forall x)P(x) \rightarrow q$$

$$(\forall x)(p \rightarrow Q(x)) = p \rightarrow (\forall x)Q(x)$$

$$(\exists x)(p \rightarrow Q(x)) = p \rightarrow (\exists x)Q(x)$$

分配律

We use laws to convert formulas to PNF.
More laws in 5.1, 5.2.



Prenex Normal Form (前束范式)

$$\begin{aligned}& \neg((\forall x)(\exists y)P(a, x, y) \rightarrow (\exists x)(\neg(\forall y)Q(y, b) \rightarrow R(x))) \\&= \neg(\neg(\forall x)(\exists y)P(a, x, y) \vee (\exists x)(\neg\neg(\forall y)Q(y, b) \vee R(x))) \\&= (\forall x)(\exists y)P(a, x, y) \wedge (\forall x)((\exists y)\neg Q(y, b) \wedge \neg R(x)) \\&= (\forall x)((\exists y)P(a, x, y) \wedge (\exists y)\neg Q(y, b) \wedge \neg R(x)) \\&= (\forall x)((\exists y)P(a, x, y) \wedge (\exists z)\neg Q(z, b) \wedge \neg R(x)) \\&= (\forall x)(\exists y)(\exists z)(P(a, x, y) \wedge \neg Q(z, b) \wedge \neg R(x)) \\&= (\forall x)(\exists y)(\exists z)S(a, b, x, y, z)\end{aligned}$$

Skolem Normal Form (Skolem标准形)

DEFINITION

A formula is in Skolem normal form if it is in prenex normal form with only universal quantifiers.

THEOREM

Every formula A can be converted to corresponding Skolem normal form B . A is unsatisfiable iff B is unsatisfiable.

This transformation does not preserve semantics but it does preserve satisfiability.

Skolem Normal Form (Skolem标准形)

$$(\forall x)(\exists y)(\exists z)S(a, b, x, y, z)$$



y depends on x

$$(\forall x)(\exists z)S(a, b, x, f(x), z)$$



$$(\forall x)S(a, b, x, f(x), g(x))$$

z depends on x

If x doesn't depend on any variables, we replace x with a constant.



Deduction (归结推理法)

$$(\forall x)(P(x) \rightarrow Q(x)) \wedge (\forall x)(Q(x) \rightarrow R(x)) \Rightarrow (\forall x)(P(x) \rightarrow R(x))$$

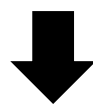
$$(\forall x)(P(x) \rightarrow Q(x)) \wedge (\forall x)(Q(x) \rightarrow R(x)) \wedge \neg(\forall x)(P(x) \rightarrow R(x))$$

$$(\forall x)(P(x) \rightarrow Q(x)) \wedge (\forall x)(Q(x) \rightarrow R(x)) \wedge (\exists x)(P(x) \wedge \neg R(x))$$

$$(\forall x)(P(x) \rightarrow Q(x)) \wedge (\forall x)(Q(x) \rightarrow R(x)) \wedge (P(a) \wedge \neg R(a))$$

$$(\forall x)((\neg P(x) \vee Q(x)) \wedge (\neg Q(x) \vee R(x)) \wedge (P(a) \wedge \neg R(a)))$$

Skolem标准形



子句集

$$\neg P(x) \vee Q(x)$$

$$\neg Q(x) \vee R(x)$$

$$P(a) \quad \neg R(a)$$

Deduction (归结推理法)

$$\neg P(x) \vee Q(x) \quad \neg Q(x) \vee R(x) \quad P(a) \quad \neg R(a)$$

$$(1) \neg P(x) \vee Q(x)$$

$$(2) \neg Q(x) \vee R(x)$$

$$(3) P(a)$$

$$(4) \neg R(a)$$

$$(5) Q(a)$$

resolution on (1)(3)

$$(6) R(a)$$

resolution on (2)(5)

$$(7) \square$$

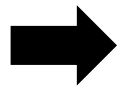
resolution on (4)(6)

Eager SMT Techniques

- Eager SMT solvers require to encode problems to SAT:
 - Represent the state space
 - Convert the problem into the new representation
- Defects: Complex encoding; producing very large formulas;

$$y = f(z) \wedge x = f(f(z)) \wedge \neg(x = f(y))$$

*x, y, z are
bool type*

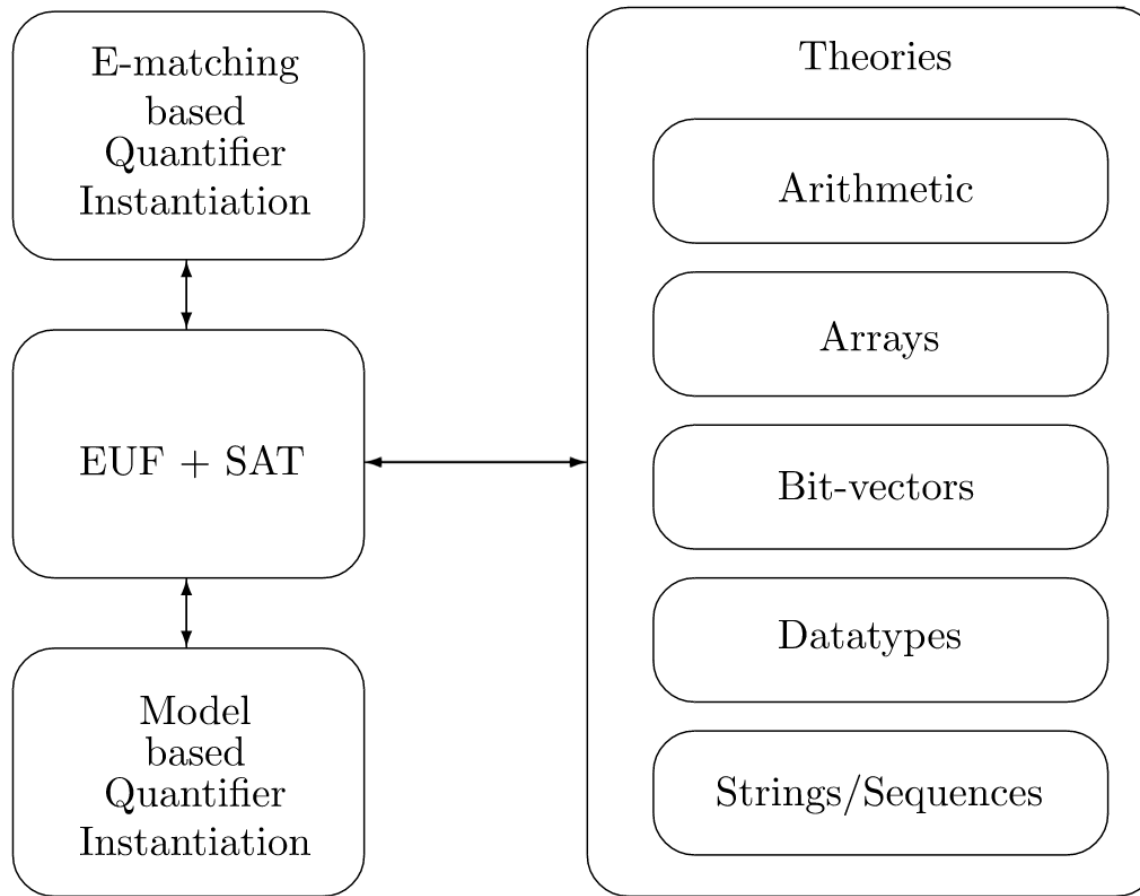


$$\begin{aligned} &((y \leftrightarrow z) \rightarrow (f_y \leftrightarrow f_z)) \wedge \\ &((y \leftrightarrow f_z) \rightarrow (f_y \leftrightarrow f_{f_z})) \wedge \\ &((z \leftrightarrow f_z) \rightarrow (f_z \leftrightarrow f_{f_z})) \wedge \\ &(y \leftrightarrow f_z) \wedge \\ &(x \leftrightarrow f_{f_z}) \wedge \\ &\neg(x \leftrightarrow f_y) \end{aligned}$$

*Property of
function f:*

$$a = b \rightarrow f(a) = f(b)$$

Lazy SMT Techniques



Integrate theory support with SAT solver

Architecture of Z3's SMT Core solver

Lazy SMT Techniques

$(x = 1 \vee x = 3) \wedge$	$(p1 \vee p2) \wedge$
$(y = 1 \vee y = 2 \vee y = 3) \wedge$	$(p3 \vee p4 \vee p5) \wedge$
$(z = 3) \wedge$	$(p6) \wedge$
$\neg(x = y) \wedge$	$\neg p7 \wedge$
$\neg(x = z) \wedge$	$\neg p8 \wedge$
$\neg(y = z)$	$\neg p9$

Step1: replace atomic formulas with fresh propositional variables.

Lazy SMT Techniques

$(x = 1 \vee x = 3) \wedge$
 $(y = 1 \vee y = 2 \vee y = 3) \wedge$
 $(z = 3) \wedge$
 $\neg(x = y) \wedge$
 $\neg(x = z) \wedge$
 $\neg(y = z)$



$(p1 \vee p2) \wedge$
 $(p3 \vee p4 \vee p5) \wedge$
 $(p6) \wedge$
 $\neg p7 \wedge$
 $\neg p8 \wedge$
 $\neg p9$



DPLL

UNSAT

If it is unsatisfiable,
the original formula is
also unsatisfiable.

Step2: use SAT solvers to
get assignments.

Lazy SMT Techniques

$$(x = 1 \vee x = 3) \wedge$$

$$(y = 1 \vee y = 2 \vee y = 3) \wedge$$

$$(z = 3) \wedge$$

$$\neg(x = y) \wedge$$

$$\neg(x = z) \wedge$$

$$\neg(y = z)$$



$$(p1 \vee p2) \wedge$$

$$(p3 \vee p4 \vee p5) \wedge$$

$$(p6) \wedge$$

$$\neg p7 \wedge$$

$$\neg p8 \wedge$$

$$\neg p9$$



DPLL

$$p1 = T, p2 = F, p3 = T, p4 = F,$$

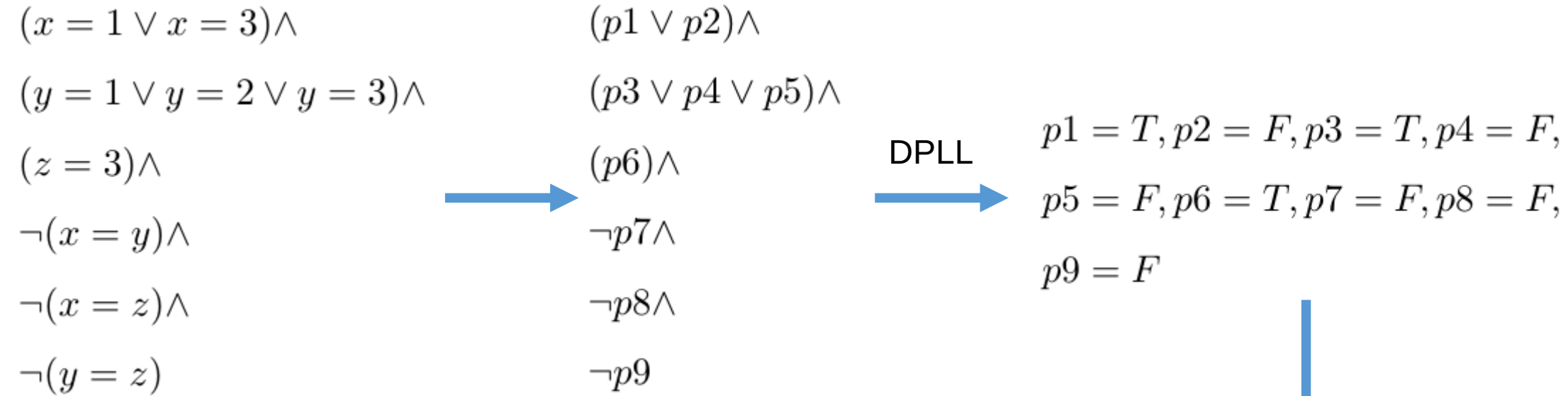
$$p5 = F, p6 = T, p7 = F, p8 = F,$$

$$p9 = F$$

DPLL will not give this result.
But to describe lazy SMT, we
suppose this is the result.

Step2: use SAT solvers to
get assignments.

Lazy SMT Techniques

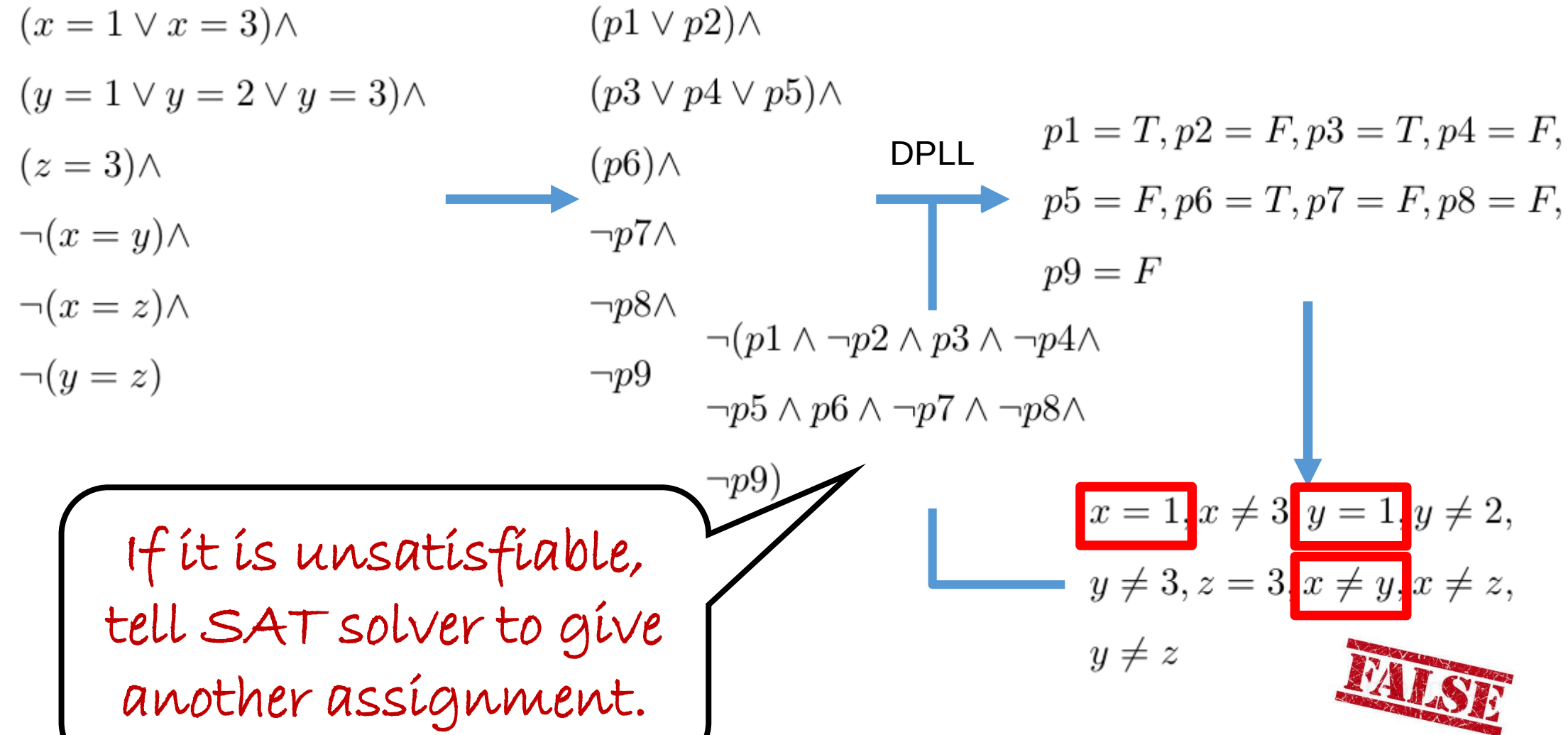


Step3: use T-solvers to check the satisfiability.

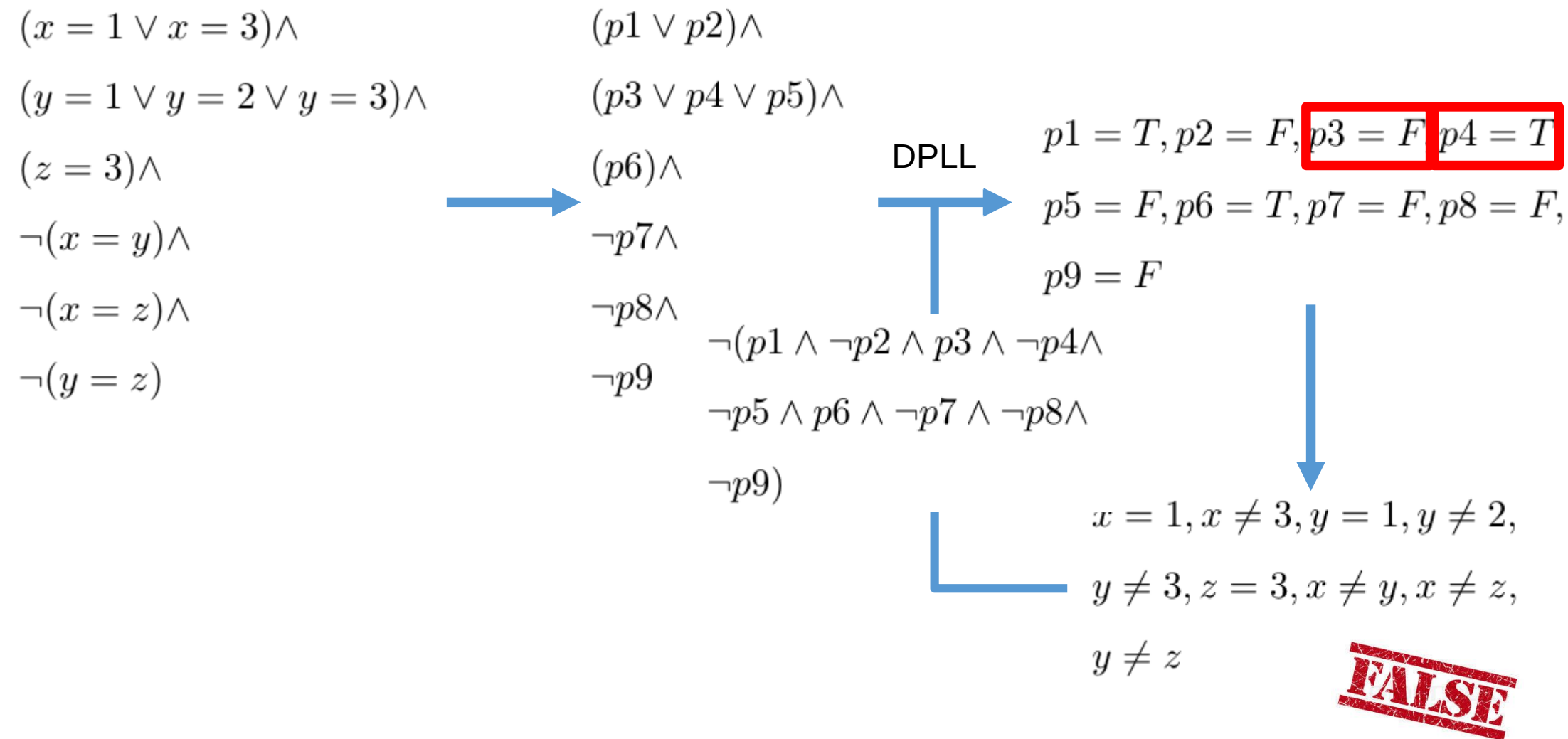
$x = 1, x \neq 3, y = 1, y \neq 2,$
 $y \neq 3, z = 3, x \neq y, x \neq z,$
 $y \neq z$

FALSE

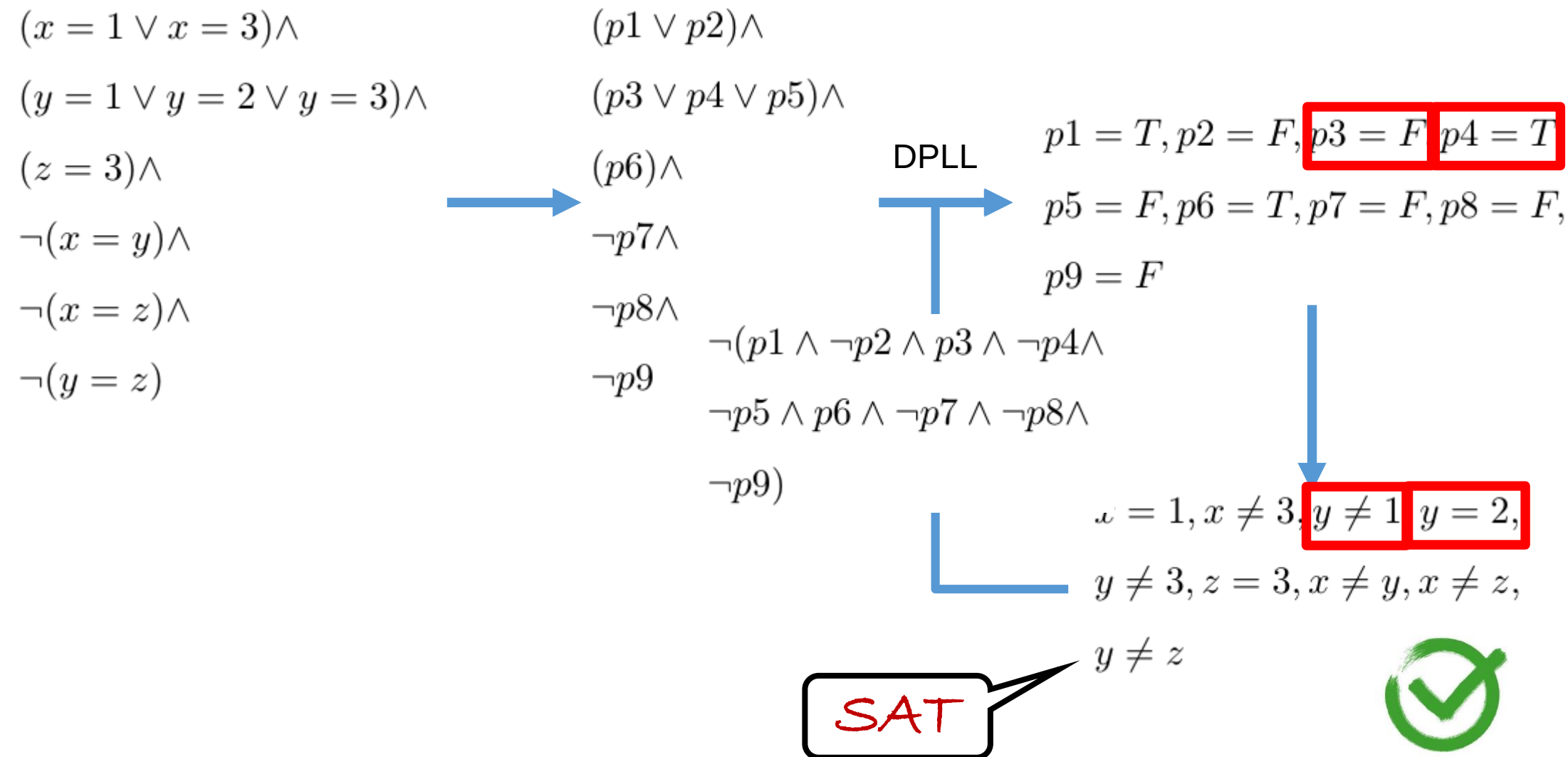
Lazy SMT Techniques



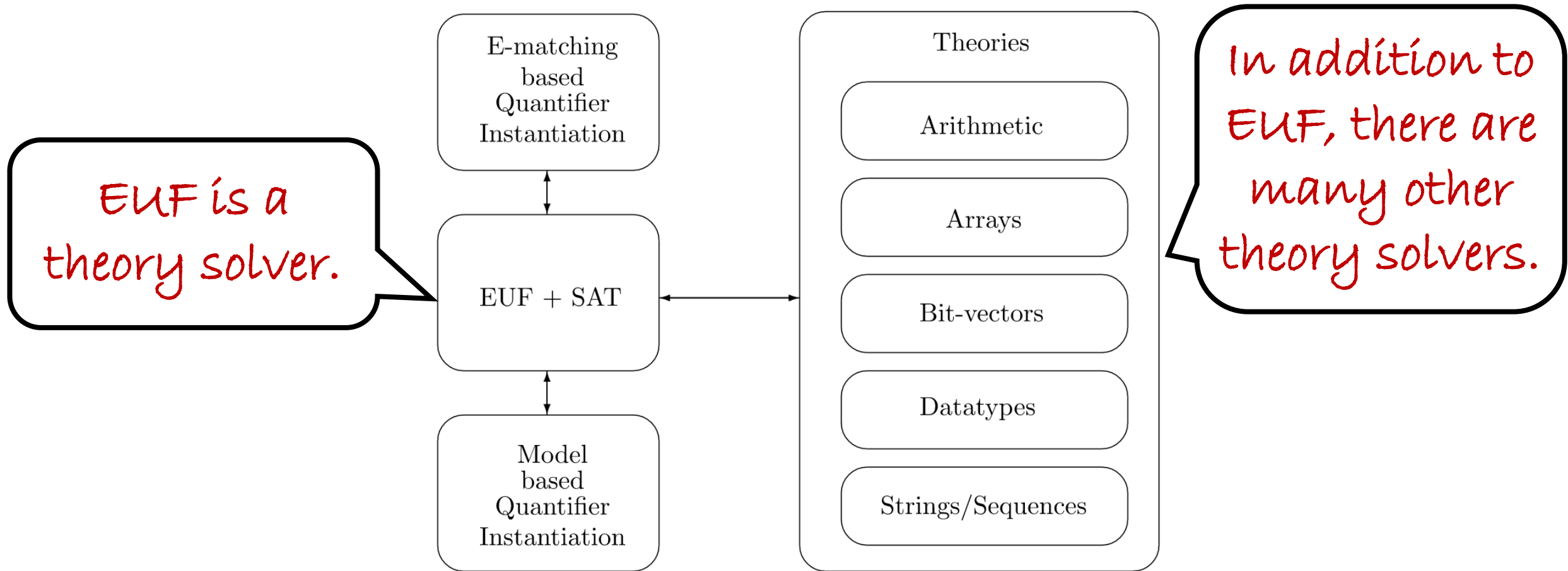
Lazy SMT Techniques



Lazy SMT Techniques



Lazy SMT Techniques



Architecture of Z3's SMT Core solver

EUF

The logic of **equality and uninterpreted function**, EUF, is a basic ingredient for (first-order) predicate logic.

While formulas are not empty:

 remove $a=b$;

 merge(a, b);

If $\text{find}(a) = \text{find}(b)$ for some $a \neq b$

 return false;

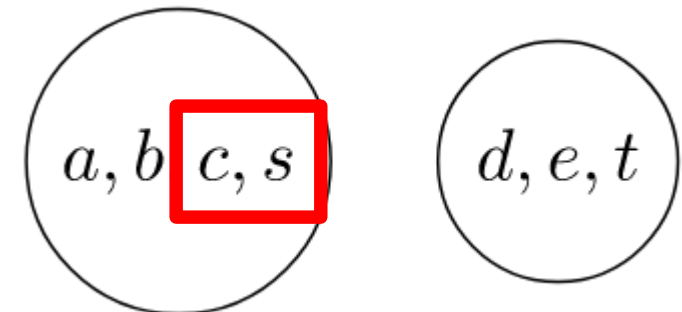
else

 return true;

$$a = b, b = c, d = e$$

$$b = s, d = t, c \neq s$$

UNSAT



EU

THEOREM

Congruence rule:

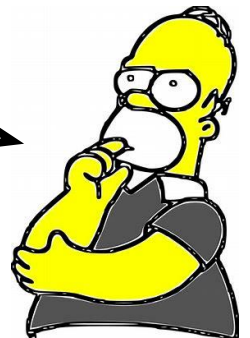
$$x_1 = y_1, \dots, x_n = y_n \Rightarrow f(x_1, \dots, x_n) = f(y_1, \dots, y_n)$$

$$a = b, b = c, d = e$$

$$b = s, d = t$$

$$f(a, g(d)) \neq f(b, g(e))$$

When there are functions,
we can construct the graph
with congruence rule.



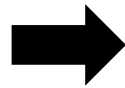
EUF

1) introduce constants that can be used as shorthands for sub-terms.

$$a = b, b = c, d = e$$

$$b = s, d = t$$

$$f(a, g(d)) \neq f(b, g(e))$$



$$a = b, b = c, d = e$$

$$b = s, d = t, v3 \neq v4$$

$$v1 := g(e), v2 := g(d)$$

$$v3 := f(a, v2), v4 := f(b, v1)$$

EUF

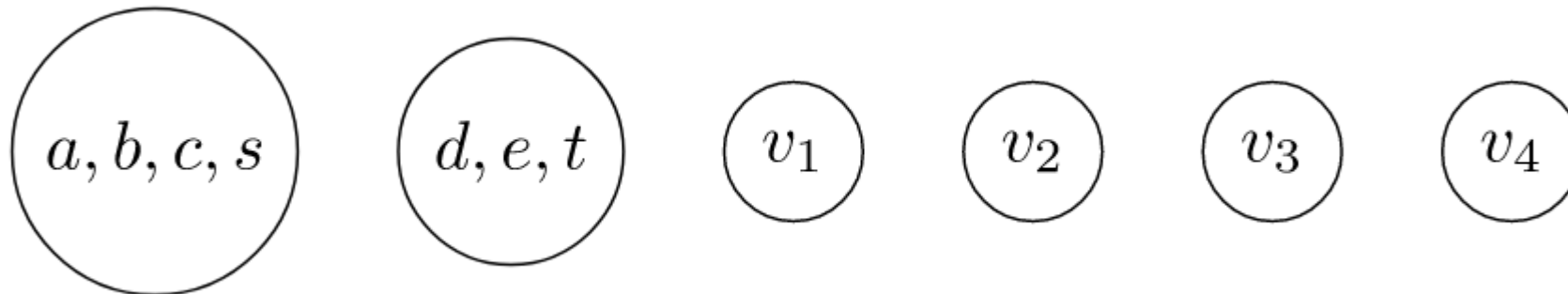
2) Working bottom-up, the congruence rule dictates how to merge groups

$$a = b, b = c, d = e$$

$$b = s, d = t, v3 \neq v4$$

$$v1 := g(e), v2 := g(d)$$

$$v3 := f(a, v2), v4 := f(b, v1)$$



EUF

2) Working bottom-up, the congruence rule dictates how to merge groups

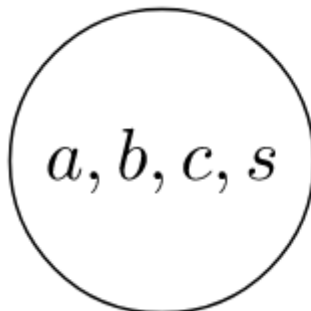
$$a = b, b = c, d = e$$

$$b = s, d = t, v3 \neq v4$$

$$v1 := g(e), v2 := g(d)$$

$$v3 := f(a, v2), v4 := f(b, v1)$$

$$e = d \Rightarrow g(e) = g(d)$$



EUF

2) Working bottom-up, the congruence rule dictates how to merge groups

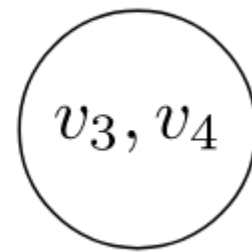
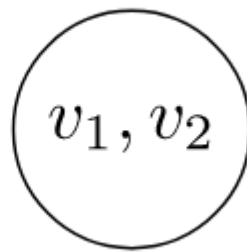
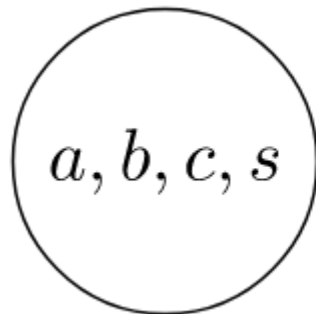
$$\begin{aligned} a = b, v2 = v1 &\Rightarrow \\ f(a, v2) = f(b, v1) \end{aligned}$$

$$a = b, b = c, d = e$$

$$b = s, d = t, v3 \neq v4$$

$$v1 := g(e), v2 := g(d)$$

$$v3 := f(a, v2), v4 := f(b, v1)$$



EUF

3) Check satisfiability like before

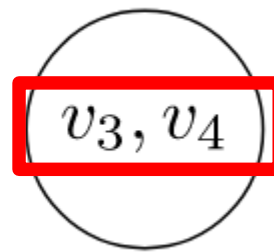
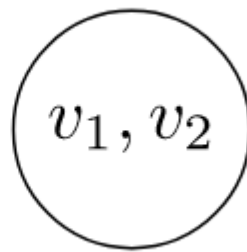
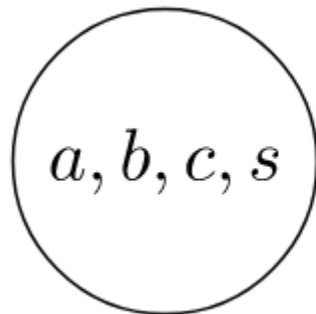
$$a = b, b = c, d = e$$

$$b = s, d = t, v3 \neq v4$$

$$v1 := g(e), v2 := g(d)$$

$$v3 := f(a, v2), v4 := f(b, v1)$$

UNSAT



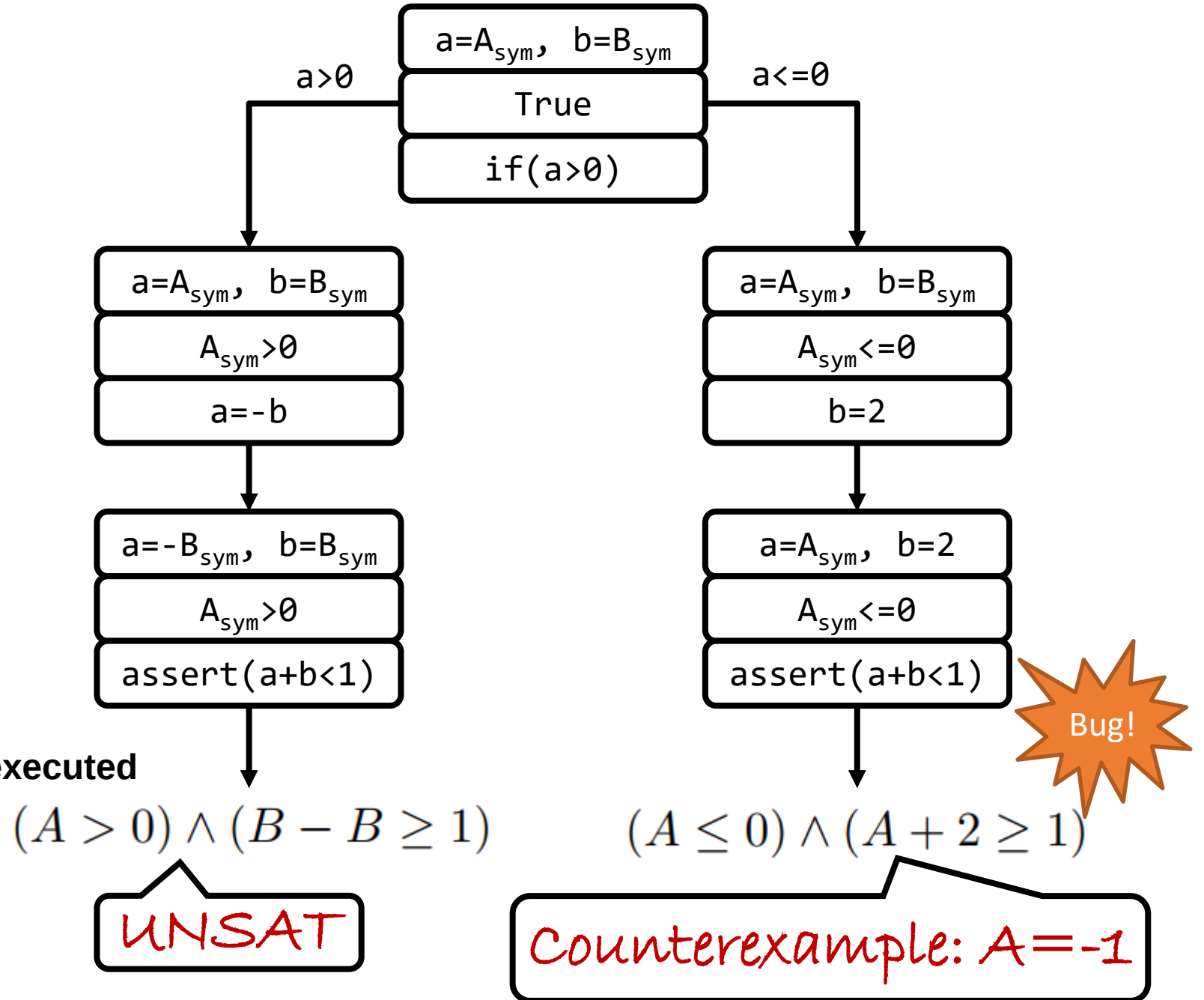
Symbolic Execution

```
void func(int a, int b) {  
    if(a > 0)  
        a = -b;  
    else  
        b = 2;  
    assert(a + b < 1);  
}
```

Path constraints: the path condition

Next code: the next instruction to be executed

Symbolic store: the symbolic value for each variable



Part I: Computer Aided Reasoning Overview

1.1 Propositional Logic

Step 1. Convert program into mathematical formula.

Step 2. Ask the computer to solve the formula.

1.2 First Order Logic

Step 1. Convert program into first order logic formula.

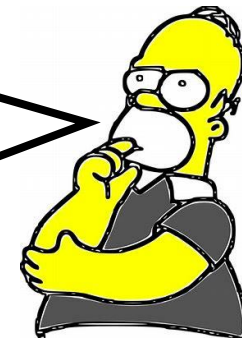
Step 2. Ask the computer to solve the formula.

1.3 Auto-active Proof

Step 1. Axiom system

Step 2. Ask the computer to check the invariants

Previous verification method cannot address unbounded loop. We introduce Hoare logic.



Defining Correctness

DEFINITION

Program state includes the value of every variable used in the program.

DEFINITION

Pre-condition describes the program state before executing the program.

DEFINITION

Post-condition describes the program state after executing the program.

pre-condition

$x = 1$

$x = x + 1;$

$x = 2$

post-condition

Hoare Triple (霍尔三元组)

DEFINITION

A program C is **partially correct** iff:

whenever C is executed in a state initially **satisfying pre-condition** and if the execution of C terminates, then the state in which C 's execution terminates **satisfies post-condition**.

pre-condition

$\{P\} C \{Q\}$

post-condition

program

Proving Correctness

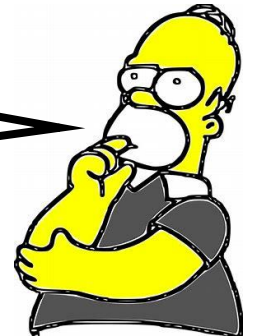
$$\{P\}a := e \{ (\exists v)(a = e[v/a] \wedge P[v/a]) \}$$

$$\{P[e/a]\}a := e \{P\}$$

$$\frac{\{P\}C1\{R\} \quad \{R\}C2\{Q\}}{\{P\}C1; C2\{Q\}}$$

$$\frac{\{P \wedge istrue(b1)\}C1\{Q\} \quad \{P \wedge isfalse(b1)\}C2\{Q\}}{\{P\}if(b1) then C1 else C2 \{Q\}}$$

A Hoare triple can be proved by these axioms and inference rules.



Loop Invariant

loop invariant

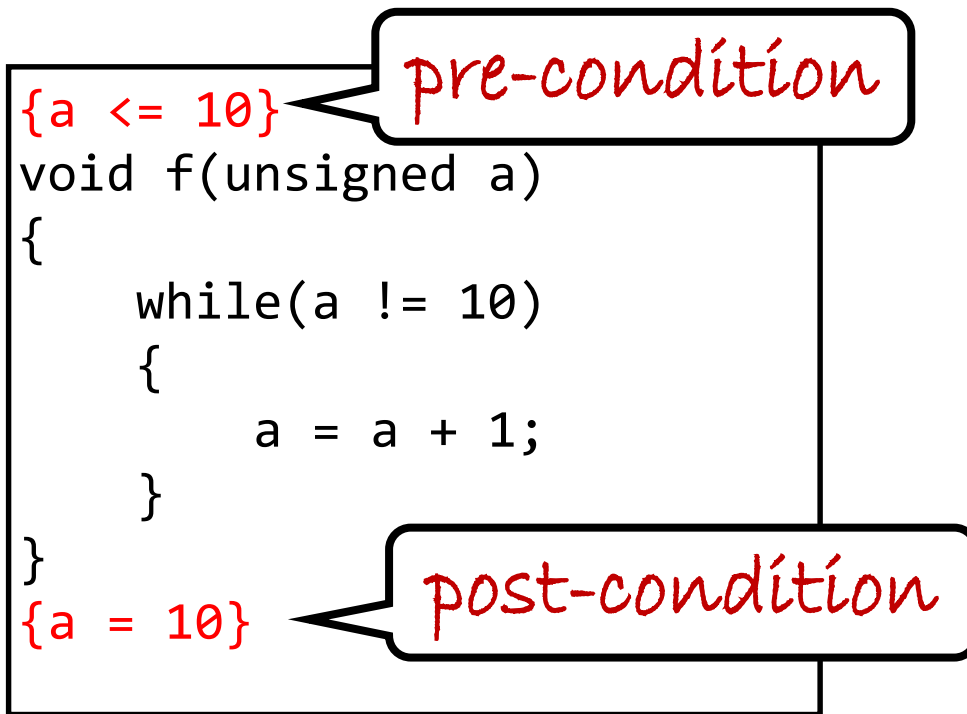
$$\frac{\{I \wedge istrue(b1)\} C \{I\}}{\{I\} \text{ while}(b1) C \{I \wedge isfalse(b1)\}}$$

The while
rule (WHP)

It says that

- if executing C once preserves the truth of I, then executing C any number of times also preserves the truth of I
- after a while loop has terminated, the test must be false

Loop Invariant



I: $a \leq 10$

loop invariant

Loop Invariant

$$\{I \wedge istrue(b1)\} C \{I\}$$

$$\{I\} \text{ while}(b1) C \{I \wedge isfalse(b1)\}$$

$$\{a \leq 10\} \text{ while}(a \neq 10) a := a + 1 \{a = 10\}$$

Proof.

$$(1) \{a + 1 \leq 10\} a := a + 1 \{a \leq 10\}$$

AS

$$(2) a \leq 10 \wedge a \neq 10 \rightarrow a + 1 \leq 10$$

predicate logic

$$(3) \{a \leq 10 \wedge a \neq 10\} a := a + 1 \{a \leq 10\}$$

SP(1)(2)

$$(4) \{a \leq 10\} \text{ while}(a \neq 10) a := a + 1 \{a \leq 10 \wedge a = 10\}$$

WHP(3)

$$(5) a \leq 10 \wedge a = 10 \rightarrow a = 10$$

predicate logic

$$(6) \{a \leq 10\} \text{ while}(a \neq 10) a := a + 1 \{a = 10\}$$

WC(4)(5)

Thanks & Questions